

Lecture 1: August 24

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This lecture's notes illustrate some uses of various \LaTeX macros. Take a look at this and imitate.

1.1 Some theorems and stuff

We now delve right into the proof.

Lemma 1.1 *This is the first lemma of the lecture.*

Proof: The proof is by induction on \dots . For fun, we throw in a figure.

Figure 1.1: A Fun Figure

This is the end of the proof, which is marked with a little box. ■

1.1.1 A few items of note

Here is an itemized list:

- this is the first item;
- this is the second item.

Here is an enumerated list:

1. this is the first item;
2. this is the second item.

Here is an exercise:

Exercise: Show that $P \neq NP$.

Here is how to define things in the proper mathematical style. Let f_k be the *AND – OR* function, defined by

$$f_k(x_1, x_2, \dots, x_{2k}) = \begin{cases} x_1 & \text{if } k = 0; \\ \text{AND}(f_{k-1}(x_1, \dots, x_{2k-1}), f_{k-1}(x_{2k-1+1}, \dots, x_{2k})) & \text{if } k \text{ is even;} \\ \text{OR}(f_{k-1}(x_1, \dots, x_{2k-1}), f_{k-1}(x_{2k-1+1}, \dots, x_{2k})) & \text{otherwise.} \end{cases}$$

Theorem 1.2 *This is the first theorem.*

Proof: This is the proof of the first theorem. We show how to write pseudo-code now.

Consider a comparison between x and y :

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if  $x$  or  $y$  or both are in  $S$  then
  answer accordingly
else
  Make the element with the larger score (say  $x$ ) win the comparison
  if  $F(x) + F(y) < \frac{n}{t-1}$  then
     $F(x) \leftarrow F(x) + F(y)$ 
     $F(y) \leftarrow 0$ 
  else
     $S \leftarrow S \cup \{x\}$ 
     $r \leftarrow r + 1$ 
  endif
endif

```

This concludes the proof. ■

1.2 Next topic

Here is a citation, just for fun [CW87].

References

- [CW87] D. COPPERSMITH and S. WINOGRAD, "Matrix multiplication via arithmetic progressions," *Proceedings of the 19th ACM Symposium on Theory of Computing*, 1987, pp. 1–6.